

Probabilistic upper bounds for the condition number of a matrix

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1 Matrix Condition number

Consider $A\mathbf{x} = \mathbf{b}$, $A \in \mathbb{R}^{n \times n}$, A nonsingular.

Problem:

- A , \mathbf{b} perturbed
- Compute $\mathbf{x} = A^{-1}\mathbf{b}$

The sensitivity of linear system:

$$\frac{\|\Delta\mathbf{x}\|}{\|\mathbf{x}\|} \leq \kappa(A) \left(\frac{\|\Delta A\|}{\|A\|} + \frac{\|\Delta\mathbf{b}\|}{\|\mathbf{b}\|} \right)$$

How to compute the condition number?

$$\kappa(A) = \|A\| \|A^{-1}\| = \frac{\sigma_{\max}(A)}{\sigma_{\min}(A)} = \sqrt{\frac{\lambda_{\max}(A^T A)}{\lambda_{\min}(A^T A)}}$$

Singular Value Decomposition

- $\approx 21n^3$ flops

Approximate the condition number

- Approximate $\sigma_{\min}(A)$ and $\sigma_{\max}(A)$

Start with a special case:

- Assume A is symmetric.
- Eigenvalues of A : $|\lambda_1| \geq \dots \geq |\lambda_n|$

Condition number of A :

$$\kappa(A) = \|A\| \|A^{-1}\| = \frac{|\lambda_1|}{|\lambda_n|}.$$

2 Bidiagonalization

For the general case:

- A nonsymmetric, $\sigma_i(A)^2 = \lambda_i(A^T A)$

$$\text{Thus } \kappa(A) = \|A\| \|A^{-1}\| = \sqrt{\frac{|\lambda_1(A^T A)|}{|\lambda_n(A^T A)|}}$$

Lanczos Bidiagonalization, procedure:

$$\mathbf{v}_0 \xrightarrow{A} \mathbf{u}_0 \xrightarrow{A^T} \mathbf{v}_1 \xrightarrow{A} \mathbf{u}_1 \xrightarrow{A^T} \dots$$

Lanczos Bidiagonalization

- $AV_k = U_k B_k$
- $A^T U_k = V_k B_k^T + \beta \mathbf{v}_{k+1} \mathbf{e}_k^T$
- $\mathbf{v}_k = p_k(A^T A) \mathbf{v}_0$
- $B_k = U_k^T A V_k$ bidiagonal
- Basis for Krylov subspace $\mathcal{K}_k(A^T A, \mathbf{v}_0)$.

HOCHSTENBACH, 2013: Probabilistic upper bound for $\|A\|$

Extended Lanczos Bidiagonalization, procedure:

$$\mathbf{v}_0 \xrightarrow{A} \mathbf{u}_0 \xrightarrow{A^T} \mathbf{v}_1 \xrightarrow{A^{-T}} \mathbf{u}_1 \xrightarrow{A^{-1}} \dots$$

$$A^TAV = VH^TH$$

$$(A^TA)^{-1}V = VKK^T$$

$$AA^TU = UHH^T$$

$$(AA^T)^{-1}U = UK^TK$$

$$\mathbf{v}_k = p_k(A^TA)\mathbf{v}_0$$

$$\mathbf{u}_k = q_k(AA^T)\mathbf{u}_0$$

- H and K are tridiagonal, $H \cdot K = I$
- p_k and q_k are Laurent polynomials
- $\mathcal{K}_{m,m}(A^TA, \mathbf{v}_0) = \text{span}\{\dots, (A^TA)^{-1}\mathbf{v}_0, \mathbf{v}_0, A^TA\mathbf{v}_0, \dots\}$.

Extended Lanczos Bidiagonalization

- Largest singular value θ_1 of H approximates $\sigma_1(A)$
- Smallest singular value θ_k of H approximates $\sigma_n(A)$
- Lower bound for condition number:

$$\frac{\theta_1}{\theta_k} \leq \kappa(A)$$

3 Probabilistic upper bound

Let $\mathbf{v}_0 = \sum_{i=1}^n \gamma_i \mathbf{y}_i$, (\mathbf{y}_i right singular vectors of A)

then

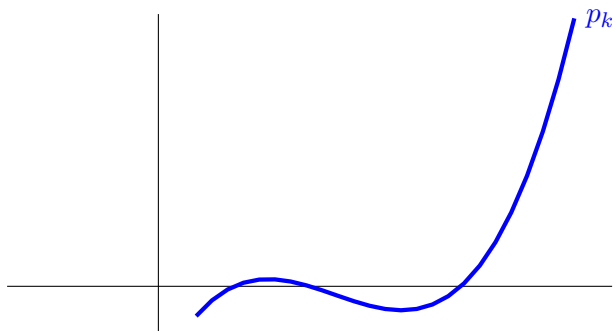
$$1 = \|\mathbf{v}_k\|^2 = \|p_k(A^T A)\mathbf{v}_0\|^2 = \sum_{i=1}^n \gamma_i^2 p_k(\sigma_i^2)^2.$$

Thus

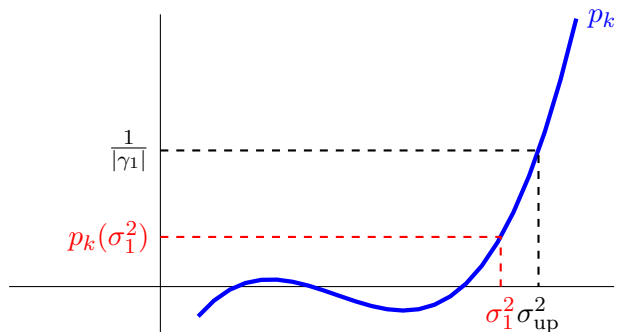
$$1 \geq \gamma_1^2 p_k(\sigma_1^2)^2,$$

and

$$\frac{1}{|\gamma_1|} \geq |p_k(\sigma_1^2)|.$$



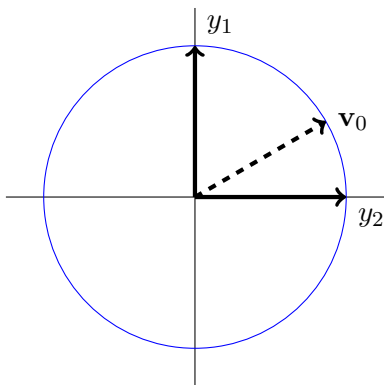
Recall: $\frac{1}{|\gamma_1|} \geq |p_k(\sigma_1^2)|$



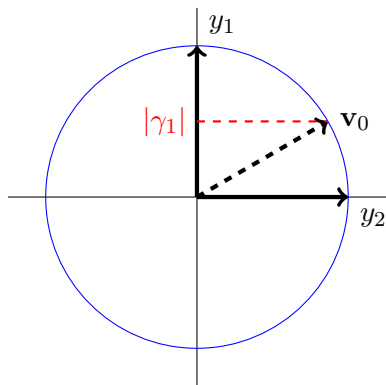
Recall: $\frac{1}{|\gamma_1|} \geq |p_k(\sigma_1^2)|$

Question: $\mathbb{P}(\frac{1}{\delta} < \frac{1}{|\gamma_1|}) = \mathbb{P}(|\gamma_1| < \delta) = \epsilon$ (ϵ is user-chosen)

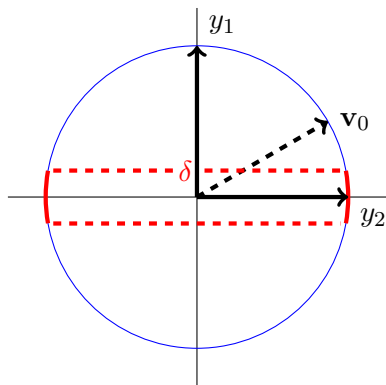
$$\mathbb{P}\left(\frac{1}{\delta} < \frac{1}{|\gamma_1|}\right) = \mathbb{P}(|\gamma_1| < \delta) = \epsilon \quad (\epsilon \text{ is user-chosen})$$



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4 Conclusions and results

Bounds for the condition number:

- lower bound: $\kappa_{\text{low}} = \frac{\theta_1}{\theta_k} \leq \kappa(A)$
- probabilistic upper bound: $\kappa(A) \leq \frac{\sigma_1^{\text{prob}}}{\sigma_n^{\text{prob}}} = \kappa_{\text{up}}$

User chosen values:

- ε , probabilistic bound holds with probability at least $1 - 2\varepsilon$
- ζ , method adaptively performs k steps such that

$$\frac{\kappa_{\text{up}}}{\kappa_{\text{low}}} \leq \zeta$$

Matrix A	Dim.	κ	κ_{low}	κ_{up}	k	CPU	LU	CPU ¹
utm5940	5940	$4.35 \cdot 10^8$	$3.98 \cdot 10^8$	$7.21 \cdot 10^8$	4	0.13	61	0.12
grcar10000	10000	$3.63 \cdot 10^0$	$3.59 \cdot 10^0$	$5.80 \cdot 10^0$	6	0.07	31	0.05
af23560	23560	$1.99 \cdot 10^4$	$1.93 \cdot 10^4$	$2.82 \cdot 10^4$	6	0.98	74	0.88
rajat16	96294	*	$5.63 \cdot 10^{12}$	$5.69 \cdot 10^{12}$	5	9.34	97	9.19
torso1	116158	*	$1.41 \cdot 10^{10}$	$1.42 \cdot 10^{10}$	3	26.8	93	28.5
dc1	116835	*	$2.39 \cdot 10^8$	$4.59 \cdot 10^8$	5	6.05	93	5.57
xenon2	157464	*	$4.29 \cdot 10^4$	$8.14 \cdot 10^4$	7	20.1	82	19.6
scircuit	170998	*	$2.40 \cdot 10^9$	$4.69 \cdot 10^9$	7	2.05	54	1.39
transient	178866	*	$1.02 \cdot 10^{11}$	$2.00 \cdot 10^{11}$	8	7.70	86	7.12
stomach	213360	*	$4.62 \cdot 10^1$	$9.02 \cdot 10^1$	6	13.8	80	13.7

- For $\zeta = 2$ (i.e. $\kappa_{\text{up}}/\kappa_{\text{low}} \leq 2$)
- $\varepsilon = 0.01$ (i.e. upper bound holds with probability at least 98%)
- CPU¹ indicates time of `condst`

Matrix A	Dim.	κ	κ_{low}	κ_{up}	k	CPU	LU
utm5940	5940	$4.35 \cdot 10^8$	$4.35 \cdot 10^8$	$4.71 \cdot 10^8$	10	0.19	42
grcar10000	10000	$3.63 \cdot 10^0$	$3.62 \cdot 10^0$	$3.97 \cdot 10^0$	13	0.13	21
af23560	23560	$1.99 \cdot 10^4$	$1.99 \cdot 10^4$	$2.12 \cdot 10^4$	9	1.14	66
rajat16	96294	*	$5.63 \cdot 10^{12}$	$5.69 \cdot 10^{12}$	5	9.34	97
torso1	116158	*	$1.41 \cdot 10^{10}$	$1.42 \cdot 10^{10}$	3	26.8	93
dc1	116835	*	$2.39 \cdot 10^8$	$2.45 \cdot 10^8$	8	6.52	91
xenon2	157464	*	$4.32 \cdot 10^4$	$4.67 \cdot 10^4$	14	23.6	70
scircuit	170998	*	$2.45 \cdot 10^9$	$2.67 \cdot 10^9$	16	3.28	33
transient	178866	*	$1.03 \cdot 10^{11}$	$1.11 \cdot 10^{11}$	21	9.47	70
stomach	213360	*	$4.82 \cdot 10^1$	$5.24 \cdot 10^1$	14	17.5	63

- For $\zeta = 1.1$ (i.e. $\kappa_{\text{up}}/\kappa_{\text{low}} \leq 1.1$)
- $\varepsilon = 0.01$ (i.e. upper bound holds with probability at least 98%)

5 References

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